

# Using Views to Implement Datalog Programs

Inverse Rules  
Duschka's Algorithm

# Inverting Rules

- ◆ Idea: “invert” the view definitions to give the global predicates definitions in terms of views and function symbols.
- ◆ Plug the globals’ definitions into the body of the query to get a direct expansion of the query into views.
- ◆ Even works when the query is a program.

# Inverting Rules --- (2)

- ◆ But the query may have function symbols in its solution, and these symbols actually have no meaning.
- ◆ We therefore need to get rid of them.
- ◆ Trick comes from Huyn -> Qian -> Duschka.

# Skolem Functions

- ◆ Logical trick for getting rid of existentially quantified variables.
- ◆ In terms of safe Datalog rules:
  - ◆ For each local (nondistinguished) variable  $X$ , pick a new function symbol  $f$  (the Skolem constant).
  - ◆ Replace  $X$  by  $f(\text{head variables})$ .

# Example

$v(X, Y) :- p(X, Z) \& p(Z, Y)$

◆ Replace  $Z$  by  $f(X, Y)$  to get:

$v(X, Y) :- p(X, f(X, Y)) \&$   
 $p(f(X, Y), Y)$

◆ Intuition: for  $v(X, Y)$  to be true, there must be some value, depending on  $X$  and  $Y$ , that makes the above body true.

# HQD Rule Inversion

- ◆ Replace a Skolemized view definition by rules with:
  1. A subgoal as the head, and
  2. The view itself as the only subgoal of the body.

# Example

```
v(X, Y) :- p(X, f(X, Y)) &  
          p(f(X, Y), Y)
```

becomes:

```
p(X, f(X, Y)) :- v(X, Y)  
p(f(X, Y), Y) :- v(X, Y)
```

# Running Example: Maternal Ancestors

## ◆ Global predicates:

- ◆  $m(X, Y)$  = “ $Y$  is the mother of  $X$ .”
- ◆  $f(X, Y)$  = “ $Y$  is the father of  $X$ .”

## ◆ manc rules:

r1:  $\text{manc}(X, Y) :- m(X, Y)$

r2:  $\text{manc}(X, Y) :- f(X, Z) \& \text{manc}(Z, Y)$

r3:  $\text{manc}(X, Y) :- m(X, Z) \& \text{manc}(Z, Y)$

# Example --- Continued

## ◆ The views:

v1(X, Y) :- f(X, Z) & m(Z, Y)

v2(X, Y) :- m(X, Y)

## ◆ Inverse rules:

r4: f(X, g(X, Y)) :- v1(X, Y)

r5: m(g(X, Y), Y) :- v1(X, Y)

r6: m(X, Y) :- v2(X, Y)

# Evaluating the Rules

- ◆ Treat views as EDB.
- ◆ Apply seminaïve evaluation to query (Datalog program).
- ◆ In general, function symbols -> no convergence.

# Evaluating the Rules

- ◆ But here, all function symbols are in the heads of global predicates.
- ◆ These are like EDB as far as the query is concerned, so no nested function symbols occur.
- ◆ One level of function symbols assures convergence.

# Example

r1: manc(X, Y) :- m(X, Y)

r2: manc(X, Y) :- f(X, Z) & manc(Z, Y)

r3: manc(X, Y) :- m(X, Z) & manc(Z, Y)

r4: f(X, g(X, Y)) :- v1(X, Y)

r5: m(g(X, Y), Y) :- v1(X, Y)

r6: m(X, Y) :- v2(X, Y)

◆ Assume v1(a,b).

# Example --- (2)

r1: manc(X, Y) :- m(X, Y)

r2: manc(X, Y) :- f(X, Z) & manc(Z, Y)

r3: manc(X, Y) :- m(X, Z) & manc(Z, Y)

r4: f(a, g(a, b)) :- v1(a, b)

r5: m(g(a, b), b) :- v1(a, b)

r6: m(X, Y) :- v2(X, Y)

◆ Assume v1(a, b).

# Example --- (3)

r1:  $\text{manc}(g(a,b), b) \ :- \ m(g(a,b), b)$

r2:  $\text{manc}(X, Y) \ :- \ f(X, Z) \ \& \ \text{manc}(Z, Y)$

r3:  $\text{manc}(X, Y) \ :- \ m(X, Z) \ \& \ \text{manc}(Z, Y)$

r4:  $f(a, g(a, b)) \ :- \ v1(a, b)$

r5:  $m(g(X, Y), Y) \ :- \ v1(X, Y)$

r6:  $m(X, Y) \ :- \ v2(X, Y)$

◆ Assume  $v1(a, b)$ .

# Example --- (4)

```
r1: manc(X,Y) :- m(X,Y)
r2: manc(a,b) :- f(a,g(a,b)) &
                  manc(g(a,b),b)
r3: manc(X,Y) :- m(X,Z) & manc(Z,Y)
r4: f(X,g(X,Y)) :- v1(X,Y)
r5: m(g(X,Y),Y) :- v1(X,Y)
r6: m(X,Y) :- v2(X,Y)
```

◆ Assume  $v1(a,b)$ .

# Example --- Concluded

- ◆ Notice that given  $v1(a,b)$ , we were able to infer  $manc(a,b)$ , even though we never found out what the value of  $g(a,b)$  [the father of  $a$ ] is.

# Rule-Rewriting

- ◆ Duschka's approach moves the function symbols out of the semiaïve evaluation and into a rule-rewriting step.
- ◆ In effect, the function symbols combine with the predicates.
  - ◆ Possible only because there are never any nested function symbols.

# Necessary Technique: Unification

- ◆ We unify two atoms by finding the simplest substitution for the variables that makes them identical.
- ◆ Linear-time algorithm known.

# Example

- ◆ The unification of  $p(f(X,Y), Z)$  and  $p(A,g(B,C))$  is  $p(f(X,Y),g(B,C))$ .
- ◆ Uses  $A \rightarrow f(X,Y); Z \rightarrow g(B,C);$  identity mapping on other variables.
- ◆  $p(X,X)$  and  $p(Y,f(Y))$  have no unification.
- ◆ Neither do  $p(X)$  and  $q(X).$

# Elimination of Function Symbols

- ◆ Repeat:
  1. Take any rule with function symbol(s) in the head.
  2. Unify that head with any subgoals, of any rule, with which it unifies.
    - ◆ But first make head variables be new, unique symbols.
- ◆ Finally, replace IDB predicates + function-symbol patterns by new predicates.

# Example

r1: manc(X, Y) :- m(X, Y)

r2: manc(X, Y) :- f(X, Z) & manc(Z, Y)

r3: manc(X, Y) :- m(X, Z) & manc(Z, Y)

r4: f(X, g(X, Y)) :- v1(X, Y)

r5: m(g(X, Y), Y) :- v1(X, Y)

r6: m(X, Y) :- v2(X, Y)

Unify

# Example --- (2)

r2: manc(X, Y) :- f(X, Z) & manc(Z, Y)

r4: f(A, g(B, C)) :-

---

r7: manc(X, Y) :- f(X, g(B, C)) &  
manc(g(B, C), Y)

Important point: in the unification of  $X$  and  $A$ , any variable could be used, but it must appear in both these places.

# Example --- (3)

r1: manc(X, Y) :- m(X, Y)

r3: manc(X, Y) :- m(X, Z) & manc(Z, Y)

r5: m(g(A, B), C) :-

---

r8: manc(g(A, B), Y) :- m(g(A, B), Y)

r9: manc(g(A, B), Y) :- m(g(A, B), Z) &  
manc(Z, Y)

# Example --- (4)

- ◆ Now we have a new pattern:  
 $\text{manc}(\text{g}(\text{A}, \text{B}), \text{C})$  .
- ◆ We must unify it with  $\text{manc}$  subgoals in r2, r3, r7, and r9.
- ◆ r2 and r7 yield nothing new, but r3 and r9 do.

# Example --- (5)

r3: manc(X, Y) :- m(X, Z) & manc(Z, Y)

r9: manc(g(A, B), Y) :- m(g(A, B), Z) &  
manc(Z, Y)

manc(g(C, D), E)

---

r10: manc(X, Y) :- m(X, g(A, B)) &  
manc(g(A, B), Y)

r11: manc(g(A, B), Y) :-  
m(g(A, B), g(C, D))  
& manc(g(C, D), Y)

# Cleaning Up the Rules

1. For each IDB predicate ( $\text{manc}$  in our example) introduce a new predicate for each function-symbol pattern.
2. Replace EDB predicates ( $m$  and  $f$  in our example) by their definition in terms of views, but only if no function symbols are introduced.

# Justification for (2)

- ◆ A function symbol in a view is useless, since the source (stored) data has no tuples with function symbols.
- ◆ A function symbol in an IDB predicate has already been taken care of by expanding the rules using all function-symbol patterns we can construct.

# New IDB Predicates

- ◆ In our example, we only need `manc1` to represent the pattern `manc(g(..),..)`.
- ◆ That is: `manc1(X,Y,Z) = manc(g(X,Y),Z).`

# Example --- r1

r1: manc(X, Y) :- m(X, Y)

r5: m(g(X, Y), Y) :- v1(X, Y)

r6: m(X, Y) :- v2(X, Y)

Substitution OK; yields

manc(X, Y) :- v2(X, Y)

Illegal --- yields g in head. Note the case  
manc(g(..),..) is taken care of by r8.

# Example --- r2 and r3

```
r2: manc(X,Y) :- f(X,Z) & manc(Z,Y)
r3: manc(X,Y) :- m(X,Z) & manc(Z,Y)
r4: f(X,g(X,Y)) :- v1(X,Y)
r5: m(g(X,Y),Y) :- v1(X,Y)
r6: m(X,Y) :- v2(X,Y)
```

Illegal --- put  
function symbol  
g in manc.

OK; yields

```
manc(X,Y) :- v2(X,Z) &
manc(Z,Y)
```

# Example --- r4, r5, r6

- ◆ The inverse rules have played their role and do not appear in the final rules.

# Example --- r7

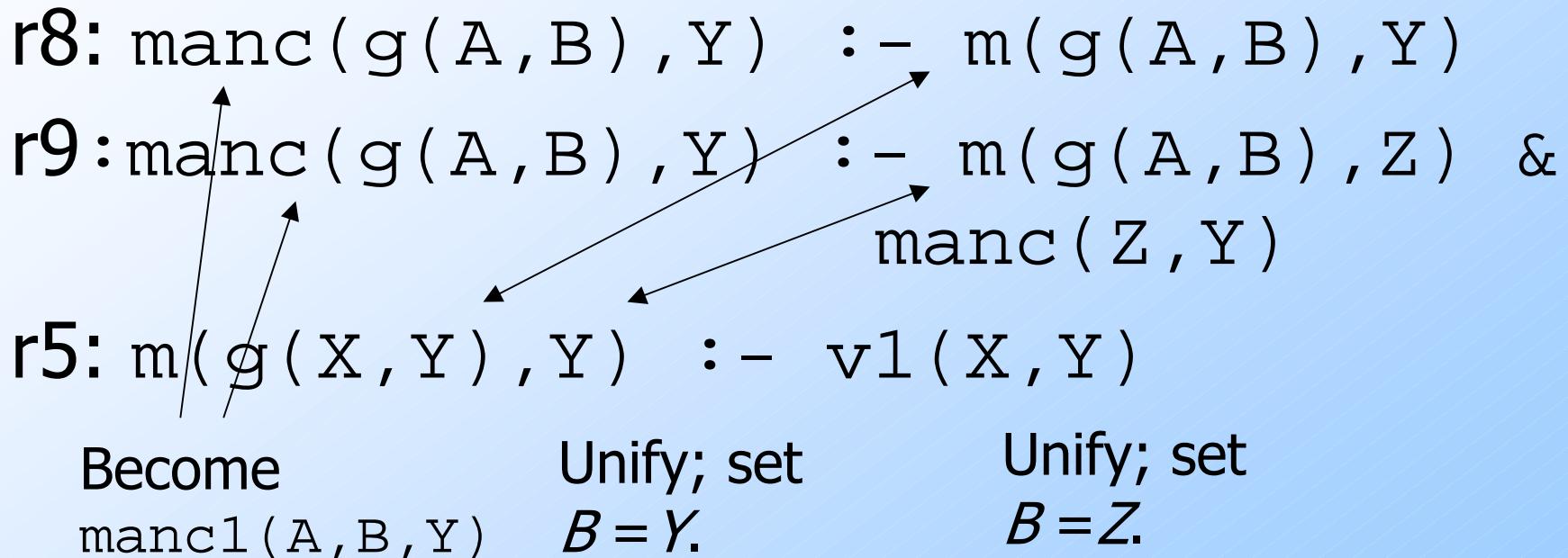
```
r7: manc(X,Y) :- f(X,g(B,C)) &  
          manc(g(B,C),Y)  
r4: f(X,g(X,Y)) :- v1(X,Y)
```

Unify. Note no  
function symbols are  
introduced, but  $X=B$ .

Replace by  
 $\text{manc1}(B,C,Y)$ .

```
r7: manc(X,Y) :- v1(X,C) &  
          manc1(X,C,Y)
```

# Example --- r8 and r9



r8:  $\text{manc1}(A, Y, Y) :- v1(A, Y)$

r9:  $\text{manc1}(A, Z, Y) :- v1(A, Z) \& \text{manc}(Z, Y)$

# Example --- r10 and r11

- ◆ No substitutions possible --- unifying  $m$  –subgoal with head of r5 or r6 introduces a function symbol into the view subgoal.

# Summary of Rules

r1: manc(X, Y) :- v2(X, Y)

r3: manc(X, Y) :- v2(X, Z) & manc(Z, Y)

r7: manc(X, Y) :- v1(X, C) &  
mancl(X, C, Y)

r8: mancl(A, Y, Y) :- v1(A, Y)

r9: mancl(A, Z, Y) :- v1(A, Z) &  
manc(Z, Y)

# Finishing Touch: Replace manc1

Substitute the bodies of r8 and r9 for the  
manc1 subgoal of r7.

r7: manc(X, Y) :- v1(X, C) &  
                              manc1(X, C, Y)

r8: manc1(A, Y, Y) :- v1(A, Y)

r9: manc1(A, Z, Y) :- v1(A, Z) &  
                              manc(Z, Y)

Becomes another  
v1(X, C).

Becomes  
v1(X, C) & manc(C, Y)

# Final Rules

- r1: manc(X, Y) :- v2(X, Y)
- r3: manc(X, Y) :- v2(X, Z) &  
                        manc(Z, Y)
- r7-8: manc(X, Y) :- v1(X, Y)
- r7-9: manc(X, Y) :- v1(X, C) &  
                        manc(C, Y)